

Lab 2: Buffer Overflows

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Buffer Overflows

One of the most common vulnerabilities in software

 Programming languages commonly associated with buffer overflows including C and C++

 Operating systems including Windows, Linux and Mac OS X are written in C or C++



How It Works

- Applications define buffers in the memory
 - Unsigned char [10]
- Applications use adjacent memory to store variables, arguments, and return address of a function.

 Buffer Overflows occurs when data written to a buffer exceeds its size.



Overflowing A Buffer

- Defining a buffer in C
 - char buf[10];

- Overflowing the buffer
 - Char buf [10] = 'x';
 - strcpy(buf, "AAAAAAAAAAAAAAAAAAAA")



Why We Care

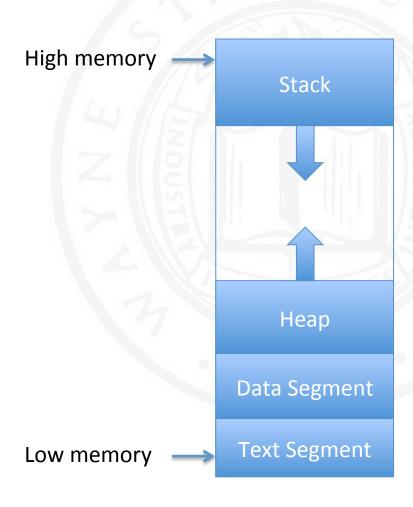
 Because adjacent memory stores program variables, parameters, and arguments

Attackers can change these values through overflowing a buffer

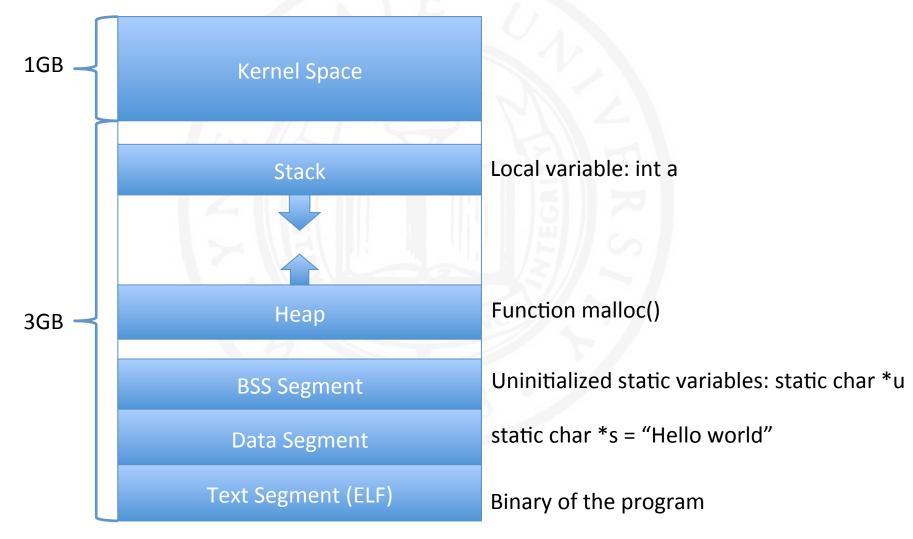
 Attackers can gain control over the program flow to execute arbitrary code



Process Memory Layout

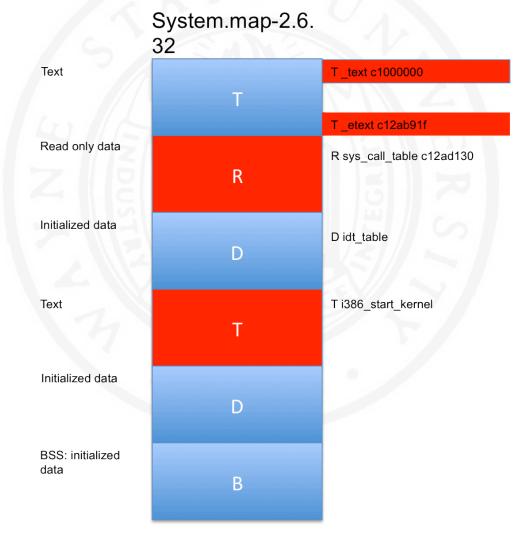


Memory Layout for 32-bit Linux





Virtual Memory Layout



Wayne State University Course: Cyber Security Practice 8



Stack Frame

 The stack contains activation frames including local variables, function parameters, and return address

Starting at the highest memory address and growing downwards

Last in first out



A Simple Program

```
Add (2,3)
                                 High memory
      int add (int a, int b)
                                                     Ret Address
           int c;
                                                         EBP
           c = 1+b;
           return c;
                                 Low memory
                                                                          ESP
```

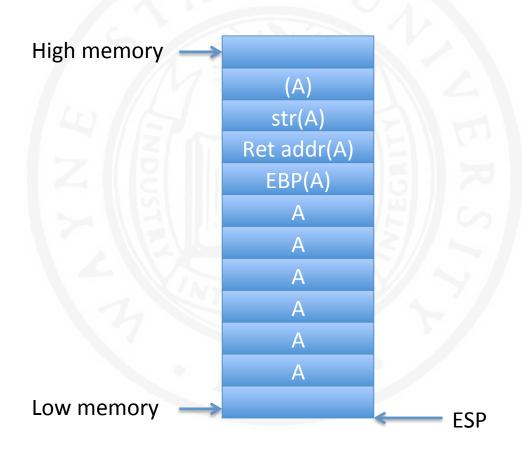


Another Program

```
int func (char * str)
   char mybuff[512];
   strcpy(myBuff, str);
   return 1;
                            Draw the Stack Frame!
int main (int argc, char ** argv)
   func (argv[1]);
   return 1;
```



Overflowing "myBuff"





Buffer Overflow Defenses

- The attack described is a classical stack smashing attack which execute the code on the stack
- It does not work today
 - NX non-executable stack. Most compilers now default to a non-executable stack. Meaning a segmentation fault occurs if running code from the stack (i.e., Data Execution Prevention - DEP)
 - Disable it with –zexecstack option
 - Check it with readelf –e <PROGRAM> | grep STACK
 - StackGuard: Cannaries
 - Disable it with –fno-stack-protector option
 - Enable it with –fstack-protector option

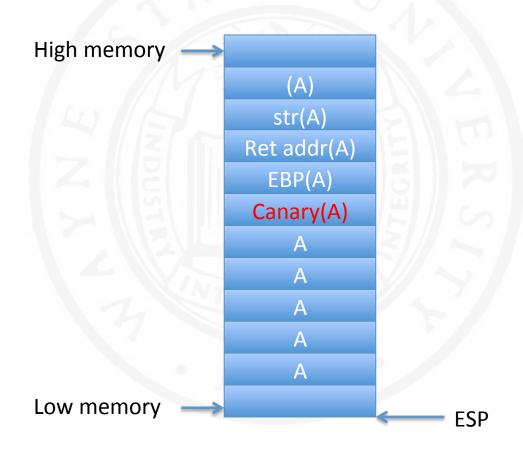


Stack Canaries

- Stack smashing attacks do two things
 - Overwrite the return address
 - Wait for algorithm to complete and call RET
- Stack Canaries: Stack Smashing Protector (SSP)
 - Placing a integer value to stack just before the return address
 - To overwrite the return address, the canary value would also be modified
 - Checking this value before the function returns



Stack Canaries (cont'd)





Bypassing NX and Canaries

- NX non-executable stack
 - Executing code in the heap
 - Data Execution Prevention (DEP)
 - Return Oriented Programming (ROP)

- Stack Canaries
 - Overwriting the Canary with the same value
 - Brute force attack (e.g., DynaGuard in ACSAC'15)



Reminders

- Lab 0
 - Turn in the class agreement
- Lab 1
 - Due today at 11:59pm
 - Late assignment policy
 - Submit it via Blackboard
- Lab 2 instructions