MobiCeal: Towards Secure and Practical Plausibly Deniable Encryption on Mobile Devices

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Outline

- Introduction
- Background
- Threat model and assumptions
- MobiCeal Design
- Performance Evaluation
- Conclusion

Introduction

• Plausibly Deniable Encryption

describes encryption techniques where the existence of an encrypted file or message is deniable in the sense that an adversary cannot prove that the plaintext data exists.

• Such denials may or may not be genuine.

Introduction cont.

- Traditional encryption systems does not work well in situations where device owner is captured by the adversary and forced to compromise the encryption key.
- Plausible Deniable Encryption (PDE) scheme can defend this kind of attack.
- Existing PDE solutions for mobile devices have few shortcomings.
 - I. Not resilient against attacks where an adversary obtains storage snapshot at different point of time.
 - II. Needs to reboot before using PDE functions.

Introduction cont.

- Existing PDE systems that work against multi snapshot adversaries also have few limitations.
 - Unsuitable for mobile environment due to performance overhead. Hive and DataLiar uses write-only oblivious RAM (ORAM) which has poor I/O performance.
 - Vulnerable to side channel attack. HIVE and DEFY do not isolate hidden and public data sufficiently.
 - Not large scale deployable. For example, DEFY is heavily dependent on a specific file system.

Background

- Full-disk encryption (FDE)
 - Full-disk encryption (FDE) is encryption at the hardware level. FDE works by automatically converting data on a hard drive into a form that cannot be understood by anyone who doesn't have the key to "undo" the conversion.

Background cont.

Plausibly Deniable Encryption

- Two types of techniques are used.
- Hidden volumes technique
 - Two encrypted volumes on the disk.
 - Public volume: encrypted using decoy key.
 - Hidden volume: encrypted using hidden key.
- Steganographic file systems
 - Hide sensitive data among regular file data.
 - Can be achieved by introducing large number of cover files or hiding data into abandoned/dummy file blocks.

Background cont.

- Logical Volume Manager (LVM)
 - LVM is a toolset that provides logical volume management capabilities in Linux.
 - LVM introduces three concepts.
 - Physical volumes (PV) : partitions or entire disk can be initialized as Physical volume
 - Volume Groups (VG) : Physical volumes are combined into volume groups
 - Logical Volumes (LV) : A volume group can be divided into logical groups.

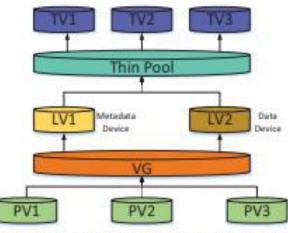
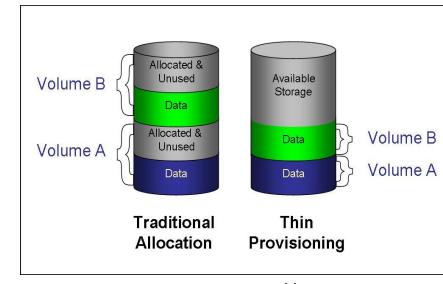


Fig. 1. LVM and thin provisioning architecture.

Background cont.

Thin Provisioning

- involves using virtualization technology to give the appearance of having more physical resources than are actually available.
- Two logical volumes are needed.
- 1. Data device: stores the data blocks.
- 2. Meta Data device: stores the space bitmap and the block mappings for thin volumes.



Source: http://wikibon.org

Threat Model and Assumptions

- Adversary can obtain multiple snapshots of the storage device.
- Adversary can have full knowledge of MobiCeal's design.
- Adversary will not capture the device when owner is using PDE functions.
- Adversary will not continue coercing owner when he is convinced encryption keys have been disclosed.
- MobiCeal code will be merged with android code stream.
- The mobile OS, kernel, bootloader, firmware, all the apps and baseband OS are malware free.

- Limitations of existing PDE in mobile devices :
 - Vulnerable to multi-snapshot adversary.
 - Attacker is dynamic, but defense is static
 - Needs dynamic defense to encounter dynamic attack.
- Limitations of dynamic defense based PDE in mobile devices :
 - Expensive in computation and I/O.
 - Designed for more dynamic attacker than our threat model.

- Uses dummy write approach.
- Perform additional artificial writes of randomness.
- Unaccountable changes in random data can be described using dummy writes.
- Number of dummy writes follows exponential distribution.
- Dummy data is created using same encryption algorithm as the hidden data.

- Creates a global bitmap to track blocks allocated for public, hidden and dummy data.
 - Sequential allocation may lead to overwrite issue.
- Current design is still vulnerable to attack when there is occasional large file written in the hidden volume.
 - Adversary may observe that public volume is followed by a large amount of randomness.

- Three types of virtual volumes.
- **Public Volume:** Public volume is encrypted using decoy key via FDE. It is used for daily operations.
- Hidden Volume: Public volume is encrypted using hidden key via FDE. Hidden volume is used to store sensitive data.
- **Dummy Volume:** Stores dummy data created by dummy writes. Purpose is to obfuscate the existence of hidden volume.

Block Allocation Strategy in Block layer

- MobiCeal uses random block allocation in the block layer.
- Sequential allocation may compromise deniability.

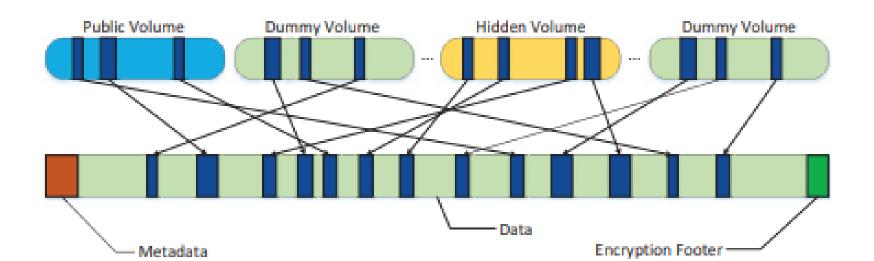
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Dv1 represents data blocks allocated for public volume and **Dv2** represents data blocks allocated to hidden volume.

- If adversary observe that seven data blocks are allocated for dummy writes between two data blocks for public voulume, he may suspect the existance of hidden volume.
- It may happen if large file is written in the hidden volume.

Storage Layout

- The entire disk is divided into three parts.
- Meta Data: stores information of the virtual volumes (global bitmap, sizes and mappings of virtual bitmap)
- Data : Stores the data block for the virtual volumes.
- Encryption Footer: Encrypted decopy key and salt are stored.

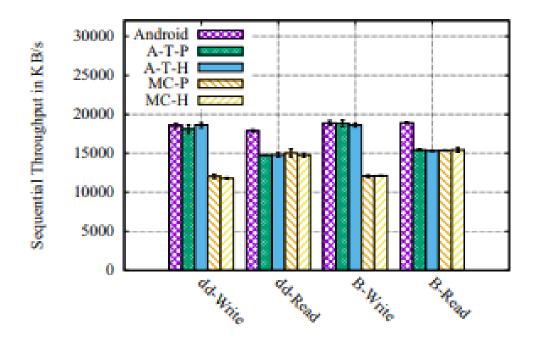


Side Channel Attacks

- Existing PDE system (e.g., HIVE and DEFY) vulnerable to side channel attacks.
- MobiCeal guards four possible leakage paths for side channel attacks.
- The public volume, logs at /devlogs, /cache and RAM are four possible leakage paths.
- To prevent side channel attacks, when hidden password is confirmed, the MobiCeal system unmounts three partitions and mounts two tmpfs RAM disks to /devlog and /cache.
- To make sure that RAM is cleared after using hidden volume, reboot is mandatory for switching to public volume.

Performance Evaluation

• Throughput Performance



Performance Evaluation

Overhead Comparison

TABLE I

OVERHEAD COMPARISON. THE VALUES OF DEFY ARE FROM THE FIGURE 6 IN [33]. TEST ENVIRONMENT: DEFY: UBUNTU 13.04, SINGLE PROCESSOR, 4GB RAM, SIMULATED FLASH DEVICE; HIVE: ARCH LINUX X86-64, I7-930, 9GB RAM, SAMSUNG 840 EVO SSD; MOBICEAL: ANDROID 4.2.2, SNAPDRAGON APQ 8064, 2GB RAM,

NEXUS 4 INTERNAL STORAGE.

	Ext4 (MB/s)	Encrypted (MB/s)	Overhead
DEFY	800	50	93.75%
HIVE	216.04	0.97	99.55%
MobiCeal	19.5	15.2	22.05%

Performance Evaluation

• Timing Measurement

TABLE II

INITIALIZATION TIME, BOOTING TIME, AND SWITCHING TIME.

	Initialization	booting time (decoy pwd)	switching time (enter hid-mod)	switching time (exit hid-mod)
Android FDE	18min23s±1s	$0.29 \pm 0.02 s$	N/A	N/A
MobiPluto	37min2s±2s	$1.36 \pm 0.02s$	68±4s	64±5s
MobiCeal	2min16s±3s	$1.68 \pm 0.04 s$	9.27±0.28s	63±6s

Conclusion

- MobiCeal is first block layer PDE scheme that is resilient to multisnapshot adversaries in mobile devices.
- It is file system friendly and supports first switching.
- Performance overhead is significantly lower comparing to other existing PDE solutions.

Thank You