## CaSE: Cache-Assisted Secure Execution on ARM Processors

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### Who am I ?

- 10 years, working on different security products data forensic, multi-level security systems
- did my undergrad @ Umass middle of no where
- did my Ph.D @ VT in DC. nice area, but I never got to go out !
- back to industry doing interesting things or not

- lastly, I am also an adjunct assistant professor at the complex network and security research laboratory (CNSR) at Virginia Tech

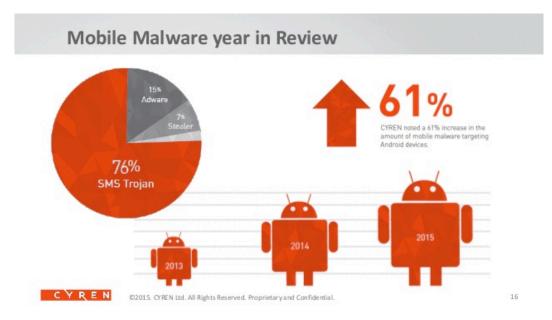
#### Talk Outline

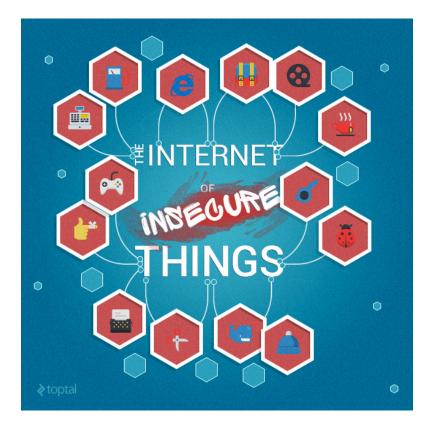
- ✓ Motivation and Background Why this work ?
- ✓ Threat Model What are we defending against ?
- ✓ CaSE: Cache-Assisted Secure Execution How does it work?
- ✓ CaSE highlight Challenges ?
- ✓ Evaluation How did we do ?
- ✓ Conclusion and future Work

#### Cyber Attacks

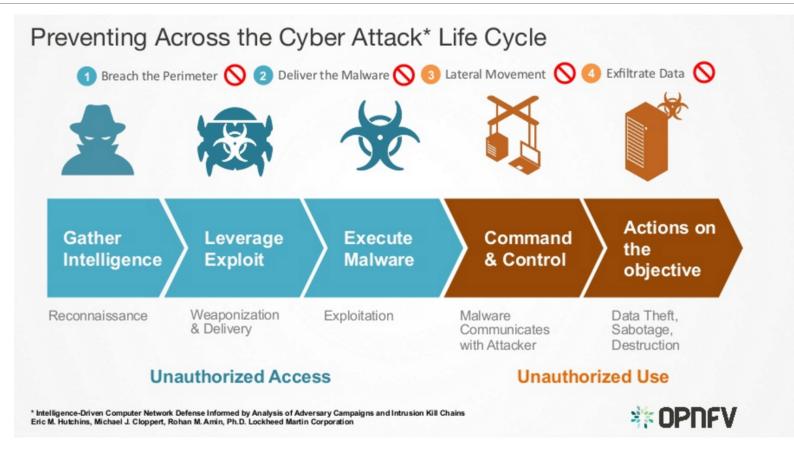


#### Threat to Mobile devices

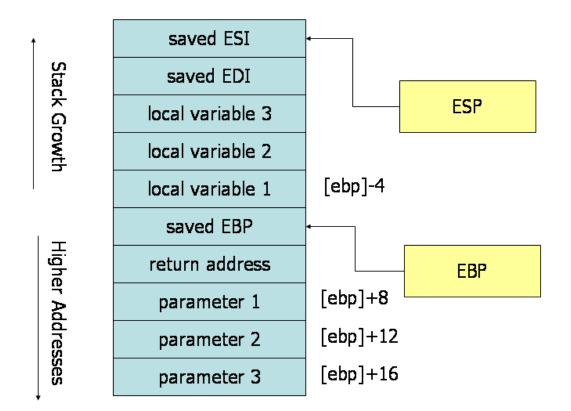




#### But how does it really work ?



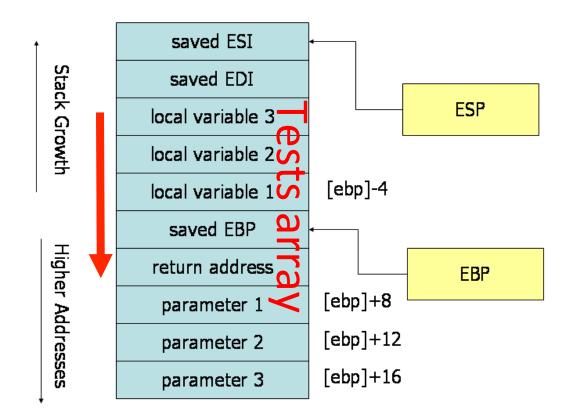
#### Buffer overflow - What is a software stack



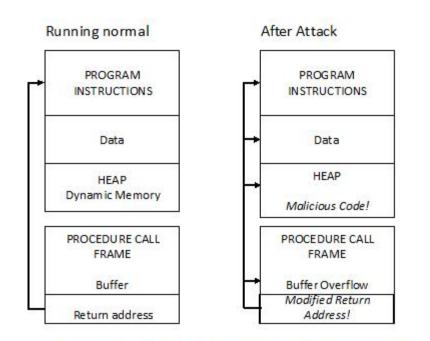
#### Software Exploits – Can you spot the bug ?

```
#include <iostream>
using namespace std;
int main(void)
{
    int tests[10];
    int test;
    int num_elems;
    cout << "How many numbers? ";
    cin >> num_elems;
    for (int i = 0; i < num_elems; i++)
        {
            cout << "Please type a number: ";
            cin >> test;
            tests[i]= test;
        }
    return 0;
}
```

#### What happened ?

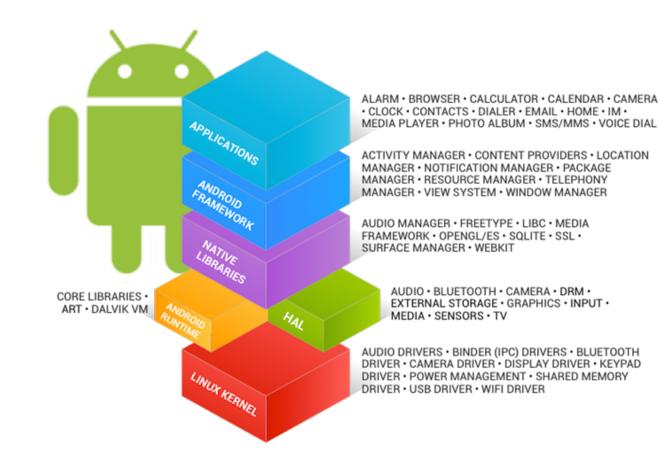


#### Before and After ③



Attacker plants code that overflows buffer and corrupts the return address. Instead of returning to the appropriate calling procedure, the modified return address returns control to malicius code, located elsewhere in process memory.

#### So are we doomed ? The best you can do ?



#### ARM TrustZone – Trusted Execution Environment (TEE)

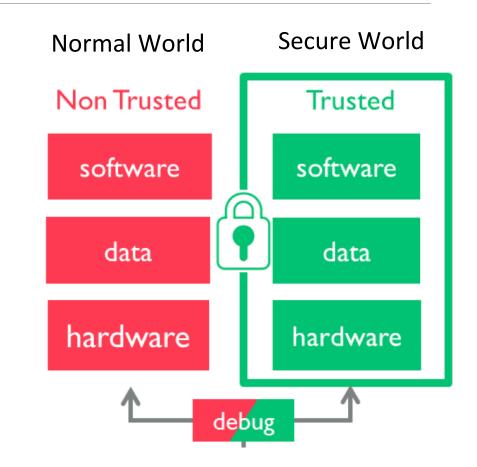
#### System Wide Protection

✓ Divides system resources into two worlds

✓ Normal World runs the content rich OS

 Secure World runs security critical services

- ✓The protection of resources includes
  - processor, memory and IO devices



#### Many Products use ARM TrustZone

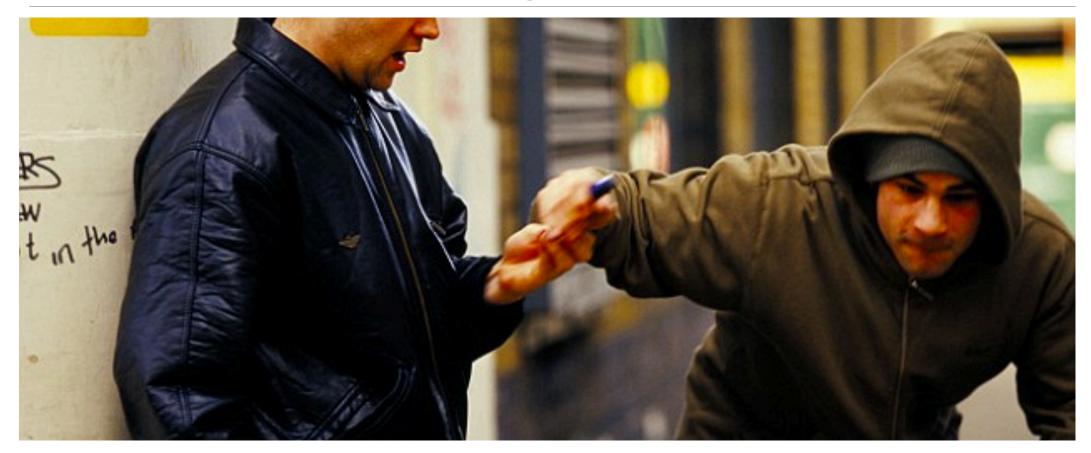
# Samsung Knax







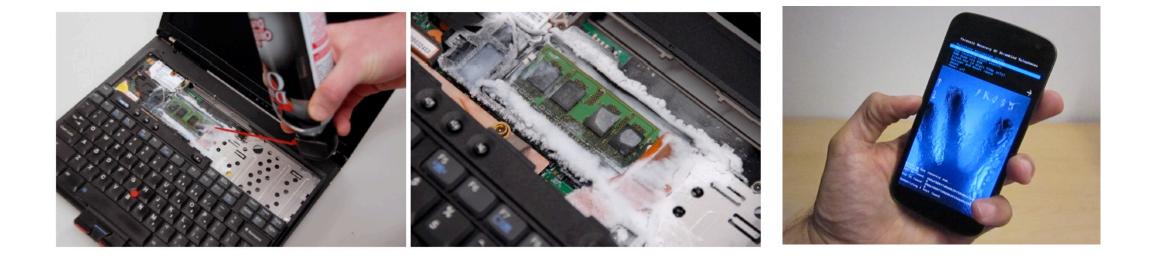
#### Smart Devices Going Mobile



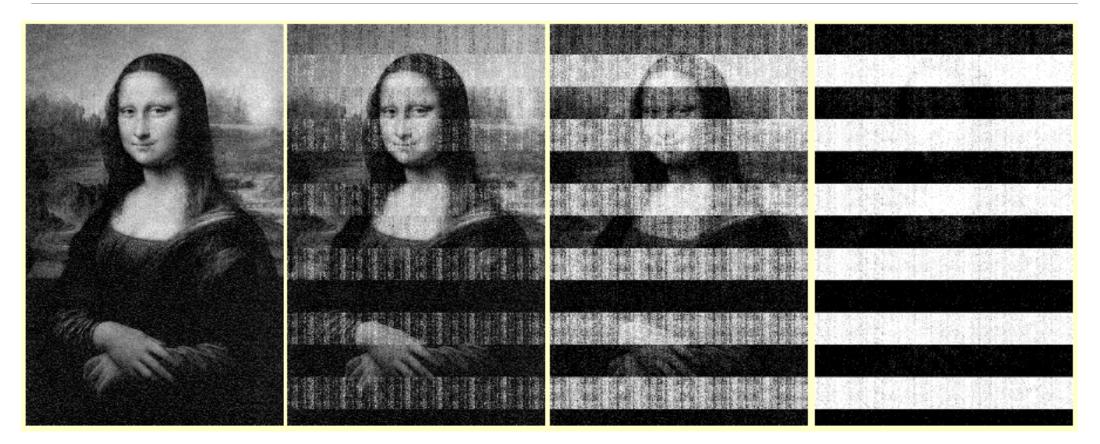
#### Physical Level Attack



#### Hardware Attacks - Cold Boot Attack



#### What can you recover ?



#### And whatever else that are in memory







#### Previous Works on Coldboot Defense

TRESOR Sec 2011 – Register-based RAM-les	AES encryption
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- Copker NDSS 2014 Cache-based RAM-less RSA encryption
- PixelVault CCS 2014 GPU based RAM-less encryption
- Sentry ASPLOS 2015 Cache-based RAM-less encryption
- Mimosa S&P 2015 Transactional-based RAM-less encryption

#### Multi-vector Adversary





#### Introducing CaSE - Goals

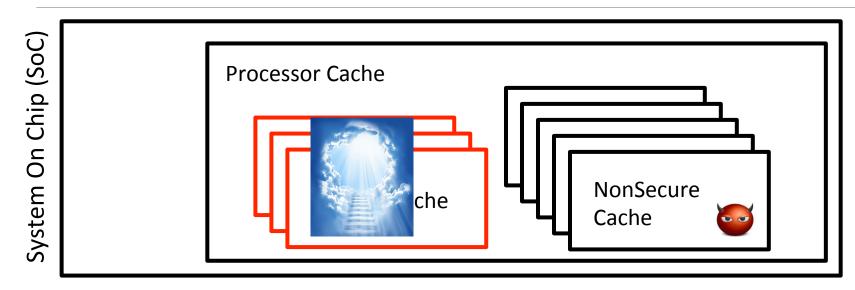
Defense against Multi-Vector adversary
 Physical memory disclosure attack – Cold boot
 Compromised rich OS

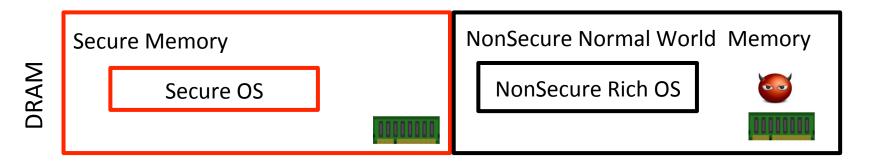
 Provide confidentiality and integrity to both the code and data of the binaries in TEE

✓ Confidentiality – Protects IP, secret code, sensitive data

✓ Integrity – Program behavior

Threat Model



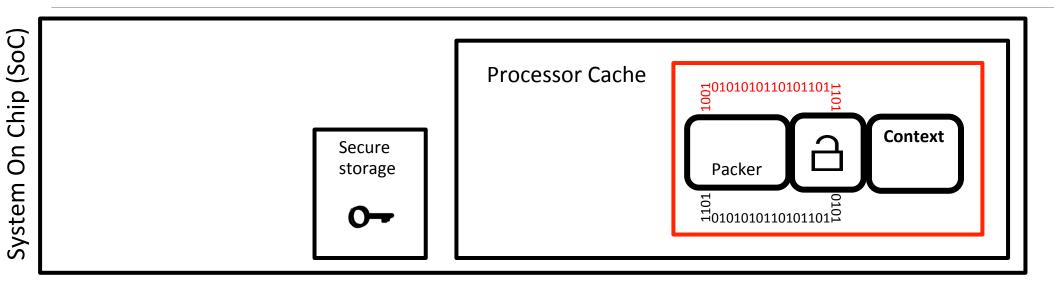


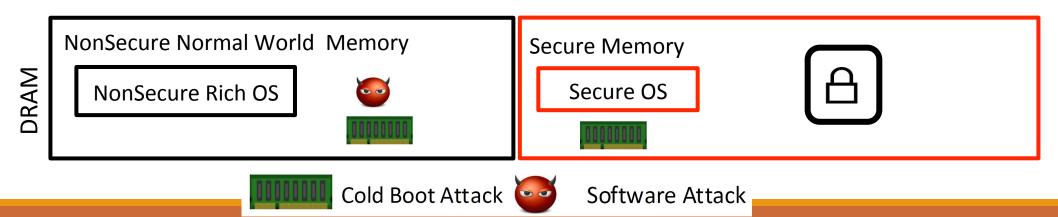




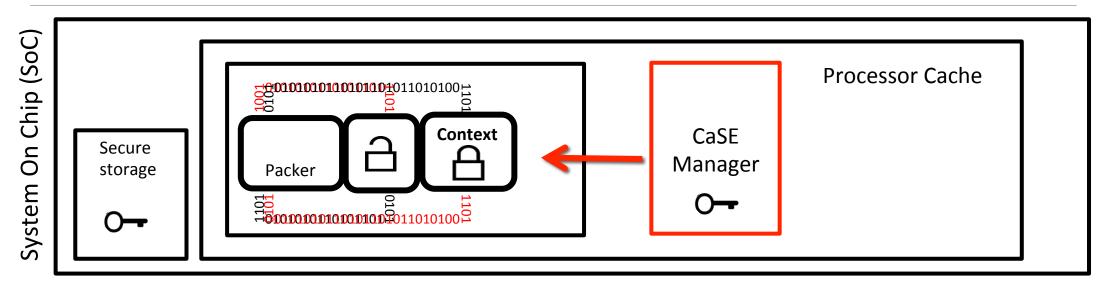
Software Attack

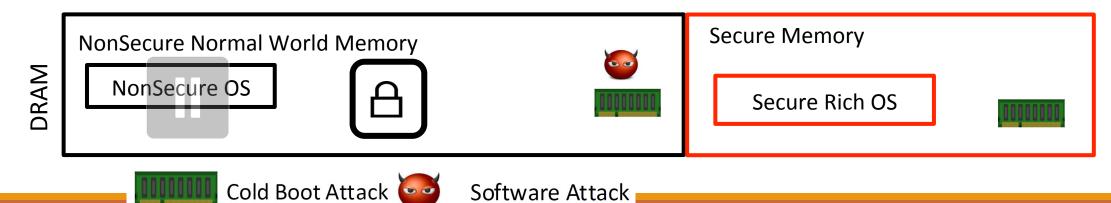
#### Case-Assisted Execution in Secure World



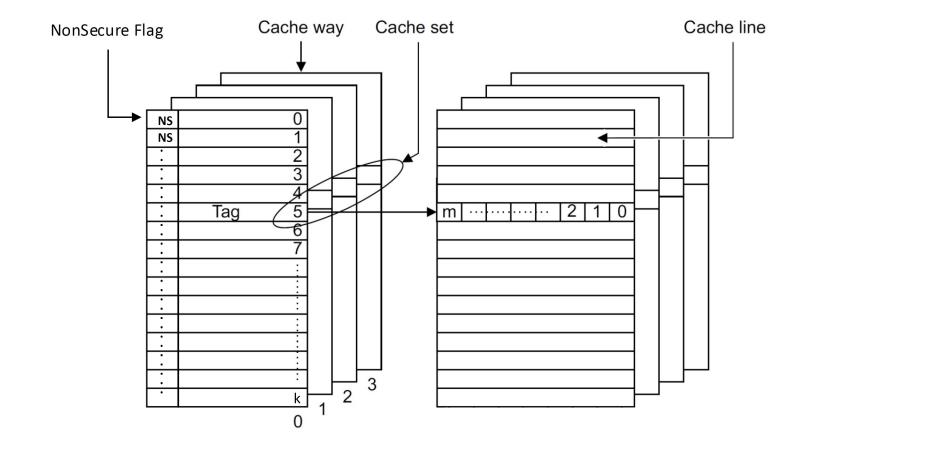


#### Case-Assisted Execution in Normal World





#### Cache Architecture Details



#### Controlling the Cache

✓ Cache Locking is available through L2 cache lockdown CP15 coprocessor

✓ The granularity of locking is per cache way

✓ On Cortex-A8, which has 8 way total 256KB L2 unified cache

#### SoC-Bound Execution – Cache Locking

```
disable_local_irq();
enableCaching(memArea);
disableCaching (loaderCode);
disableCaching (loaderStack);
invalidate_cache(virtual address of memArea);
unlockWay(wayToFill);
lockWay(allWay XOR wayToFill);
while (has more to load in memArea)
      LDR r0, [memArea + i];
lockWay(wayToFill);
unlockWay(allWay XOR wayToFill);
```

root@raspberrupi:"/ > git clone --verbose git://github.com/Hexxeh/rpi-firmware.git --depth=1 Cloning into 'rpi-firmware'... remote: Counting objects: 1673, done. remote: Compressing objects: 100% (1347/1347), done. remote: Total 1673 (delta 286), reused 1291 (delta 206) Receiving objects: 100% (1673/1673), 27.08 MiB | 306 KiB/s, done. Resolving deltas: 100% (286/286), done. [ 1461.679215] -----[ cut here ]-----[ 1461.692804] kernel BUG at drivers/tty/vt/vt.c:2838! [ 1461.706496] Internal error: Oops - BUG: 0 [#1] PREEMPT ARM Entering kdb (current=0xc5e04360, pid 1326) Oops: (null) due to oops @ 0xc0227cc8 Ω Pid: 1326, comm: agetty Tainted: G C (3.6.11 + #375) CPU: 0 PC is at con\_shutdown+0x30/0x34 LR is at queue\_release\_one\_tty+0x20/0x54 pc : [<c0227cc8>] lr : [<c02125e0>] psr: 60000013 sp : c7bedd20 ip : 00000000 fp : 0000000 r10: 00000000 r9 : 00000000 r8 : c78a41d8 r7 : 00000002 r6 : c7bec000 r5 : 00000000 r4 : c769a000 r3 : c0227c98 r2 : 00000000 r1 : 00000000 r0 : c769a000 Flags: nZCv IRQs on FIQs on Mode SVC\_32 ISA ARM Segment user Control: 00c5387d Table: 03e50008 DAC: 00000015 [<c0013a7c>] (unwind\_backtrace+0x0/0xf0) from [<c0072a80>] (kdb\_dumpregs+0x28/0x50) [<c0072a80>] (kdb\_dumpregs+0x28/0x50) from [<c0074e04>] (kdb\_main\_loop+0x3a8/0x6fc) [<c0074e04>] (kdb\_main\_loop+0x3a8/0x6fc) from [<c00774e8>] (kdb\_stub+0x154/0x380) [<c00774e8>] (kdb\_stub+0x154/0x380) from [<c006e61c>] (kgdb\_handle\_exception+0x1f8/0x668) more> \_

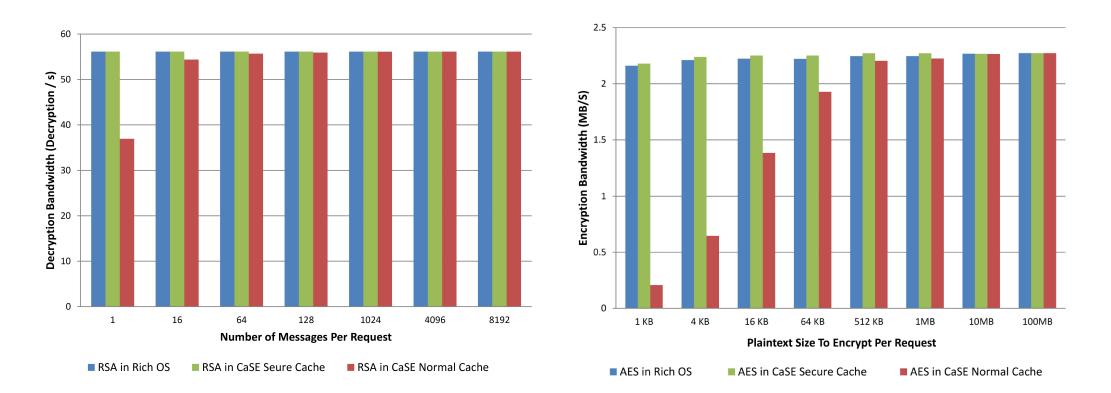
### Self Modifying Program

System On Chip (SoC)	L1 Instruction Cache		L1 Data Cache	
	L2 Unified Cache	2		

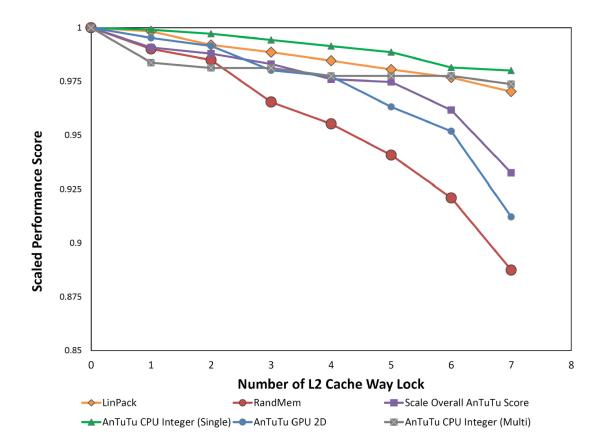
#### Evaluation Feasibility of using Cache as Memory

Application	Code+Data (KB)
AES	2.4
RSA	10
SHA1	5
CaSE Crypto Lib	17.4
Kernel Integrity Checker	6.6
CaSE Packer	2.8
Packed CaSE Crypto Lib	20.4
Packed Kernel Checker	9.5

#### Evaluation Performance Impact to the Application



#### Performance Impact to the System



#### Conclusion

- ✓A secure cache-assisted SoC-bound execution framework
  - Provide confidentiality and integrity to sensitive code and data of applications
  - ✓ Protect against both software attacks and cold boot attack.
- ✓ In the future, we would like to further study efficient method to provide OS support to the TEE.

#### What other things did I do ?

- Differential privacy in data mining ICC 11
- Reverse engineer ASUS BIOS Trusted Cloud Computing CNS 14
- Anti-memory forensic framework HIVES ASIACCS 15
- Cache-based rootkits EUROSP 16
- Case Cached-assisted security execution SP16
- Augmented reality authentication TRUSTED CCS16

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